

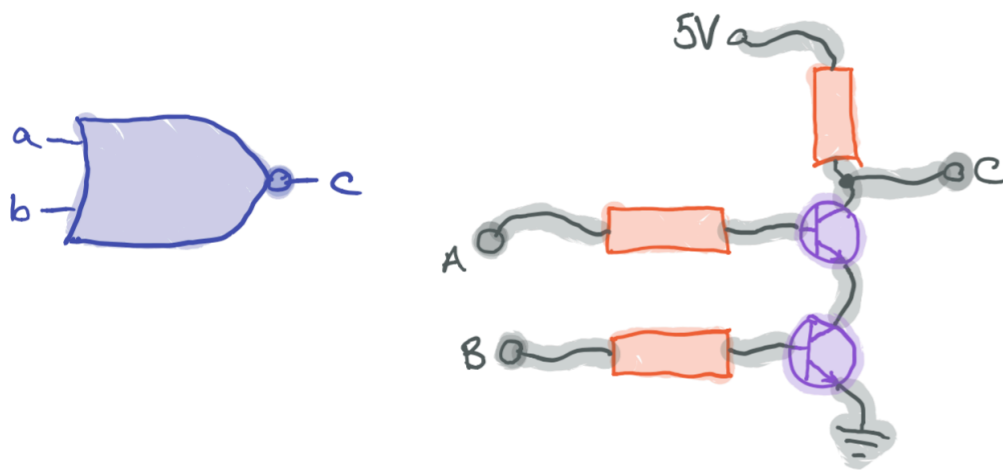
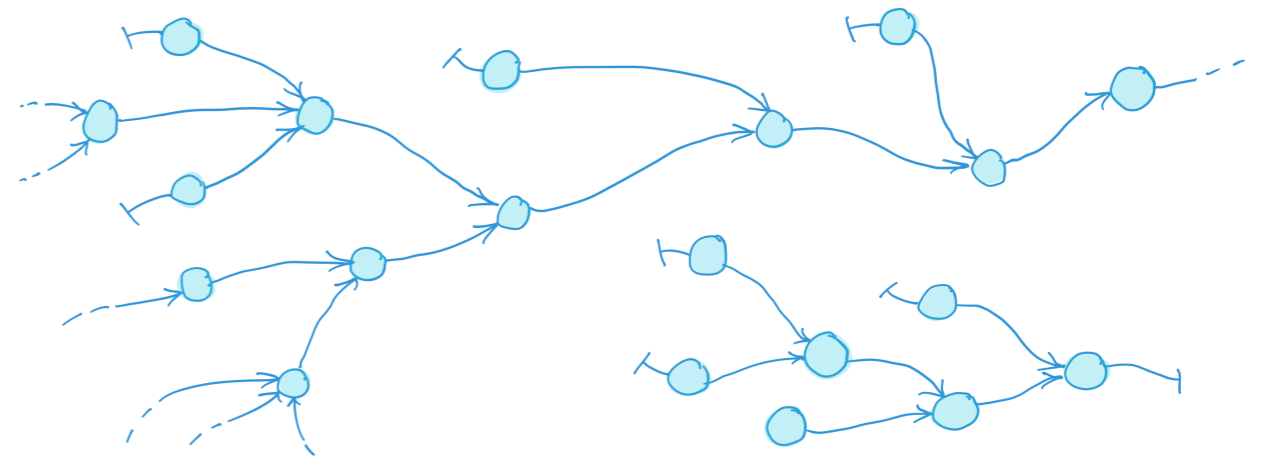
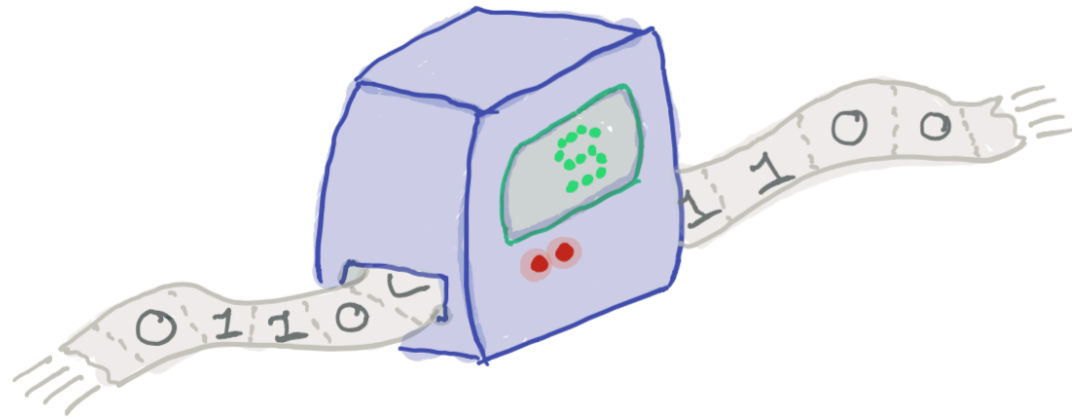
Reversible Computing:

Reuniting Computers & The Laws of Physics

William Earley

Micklem Lab · DAMTP

Irreversibility in Computing



a	b	c
0	0	0
0	1	0
1	0	0
1	1	1

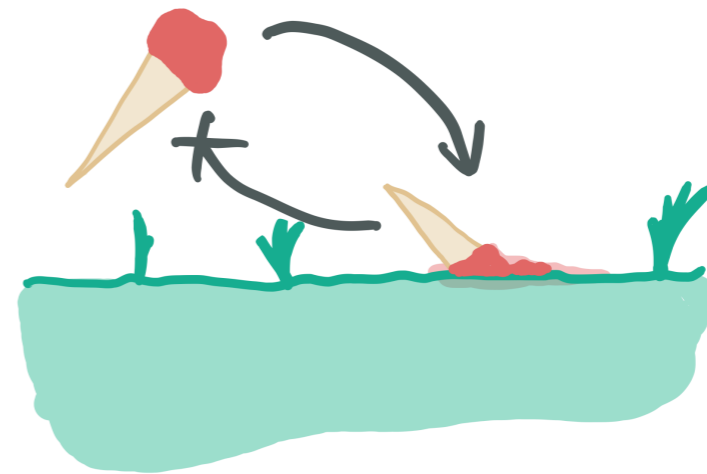
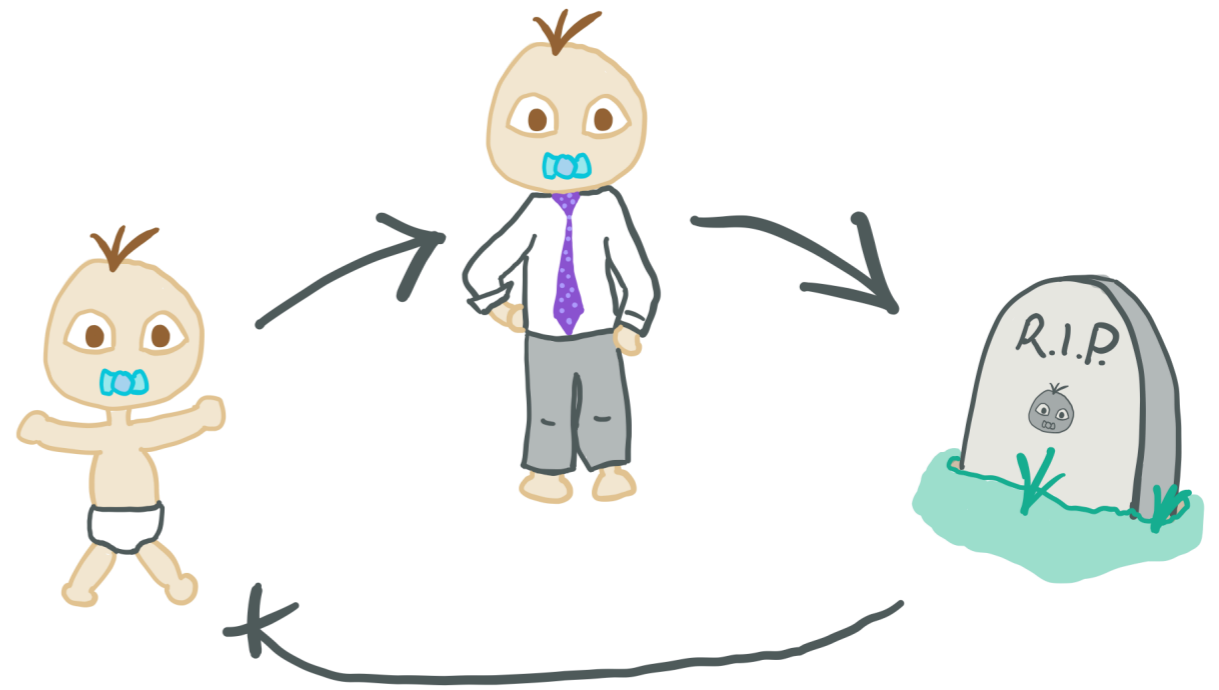
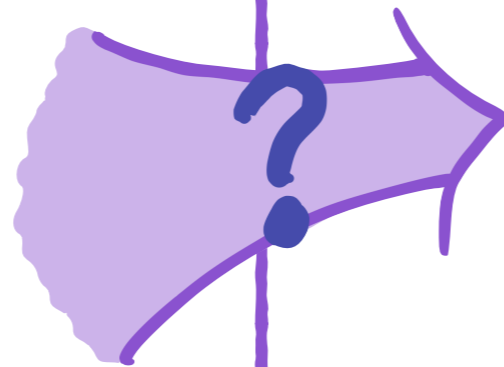
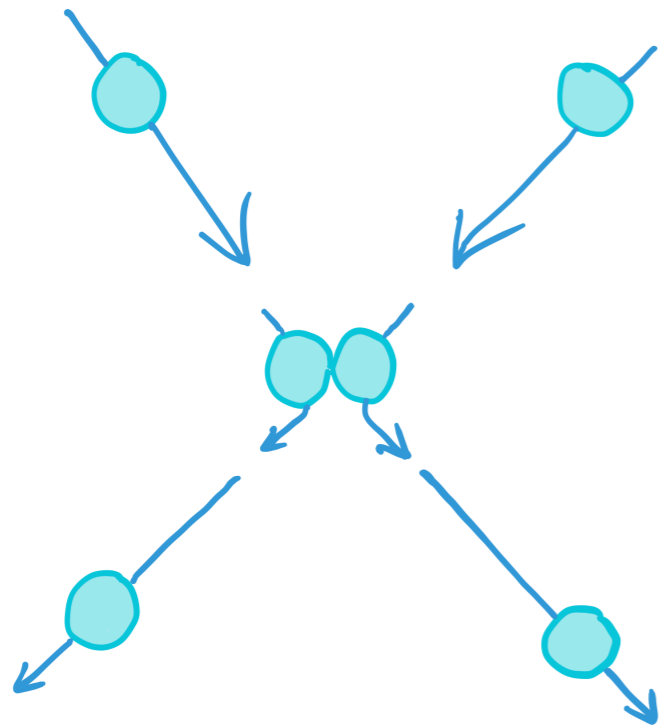
$\text{TRUE} = \lambda x. \lambda y. x$

```
>>> sum(range(10))
45
```

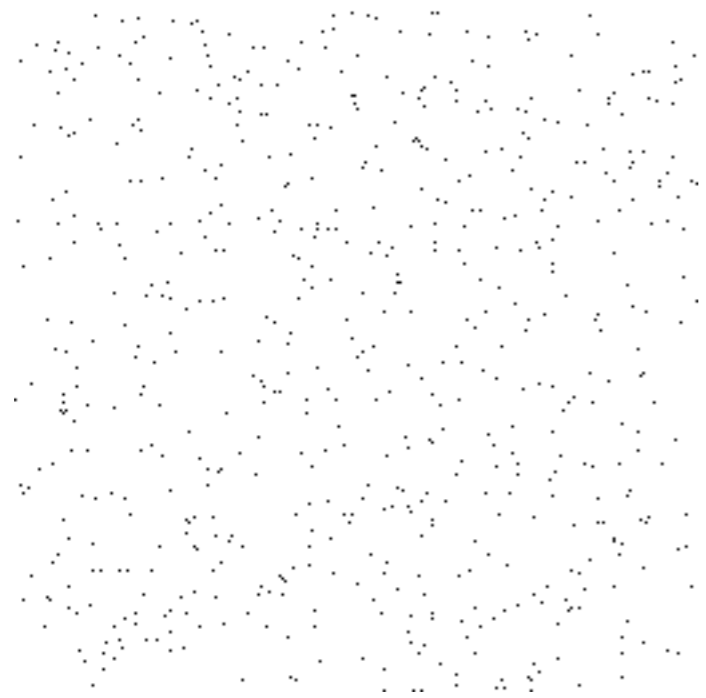
```
(curl -s 'wtr.in/Cambridge' | grep -i rain) 2>/dev/null
```

Irreversibility in Physics

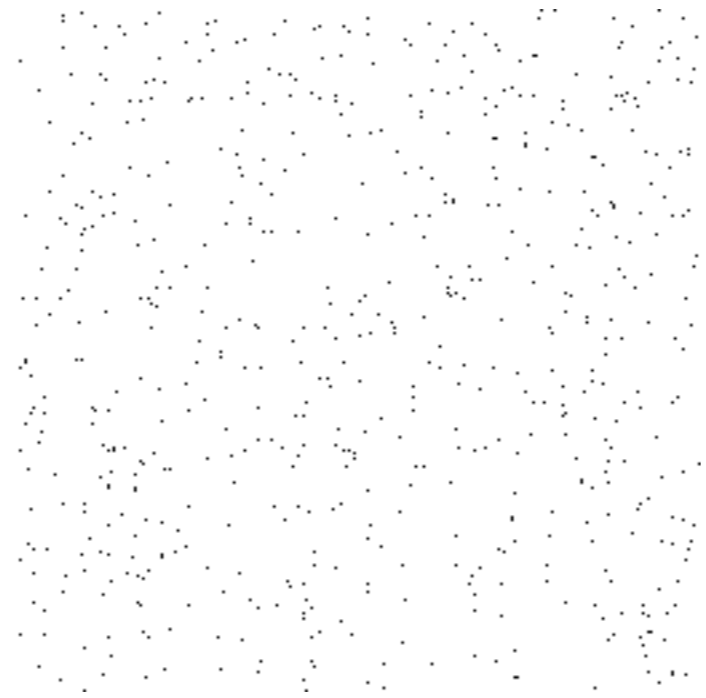
$$|\psi(t)\rangle = e^{-i\hat{H}t/\hbar} |\psi(0)\rangle$$



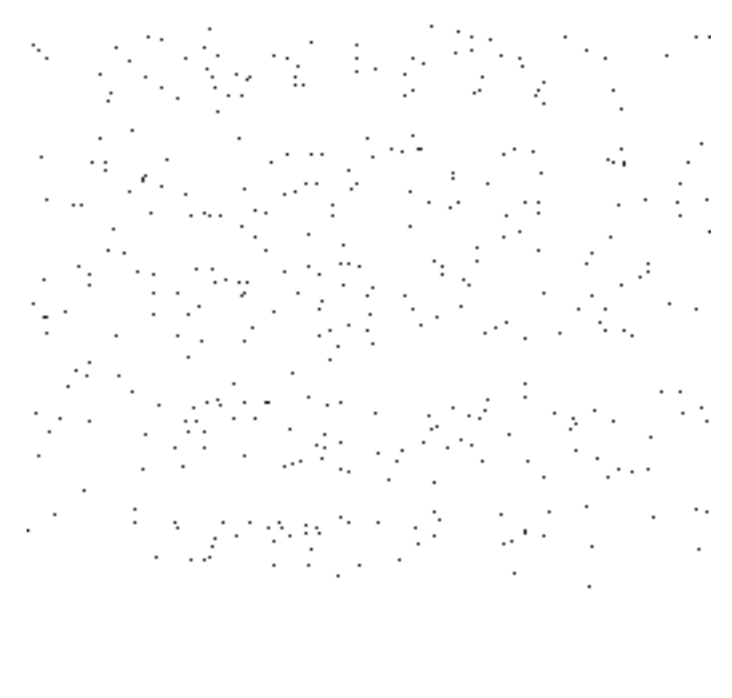
Simulating Irreversibility



[0,100]



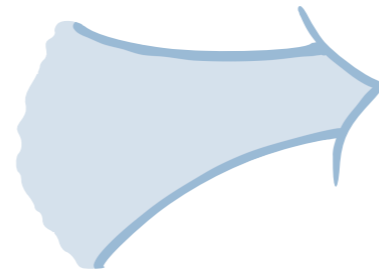
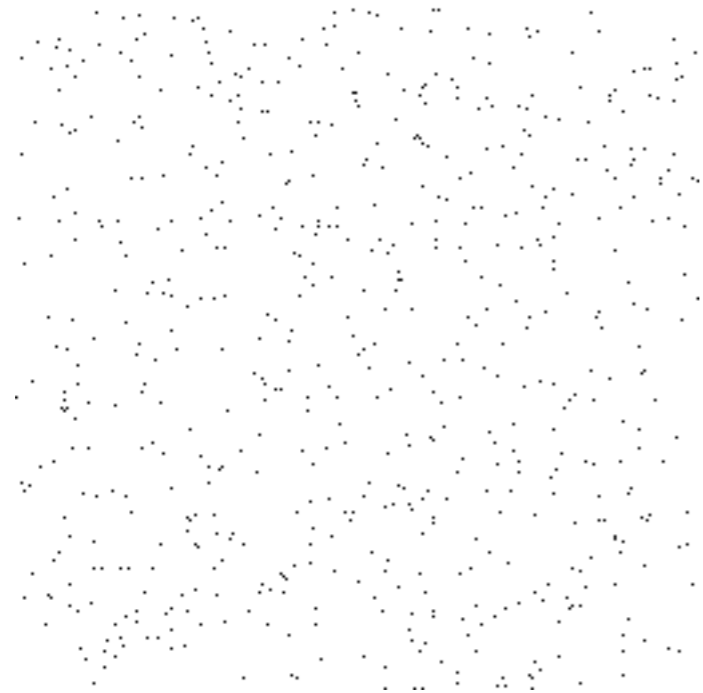
[100,0]



[-200,-300]

Asymmetry

Simulating Irreversibility

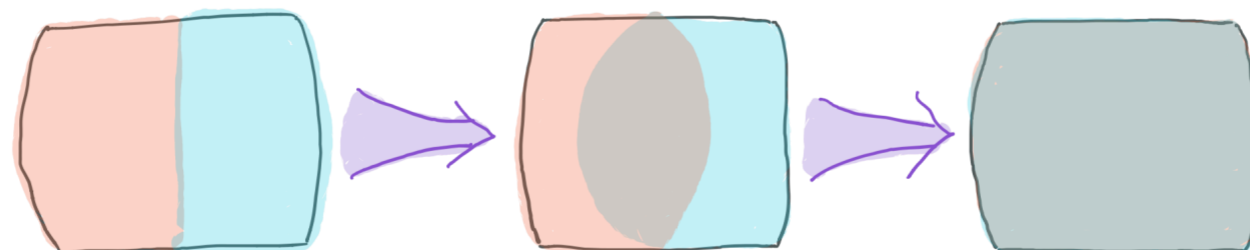
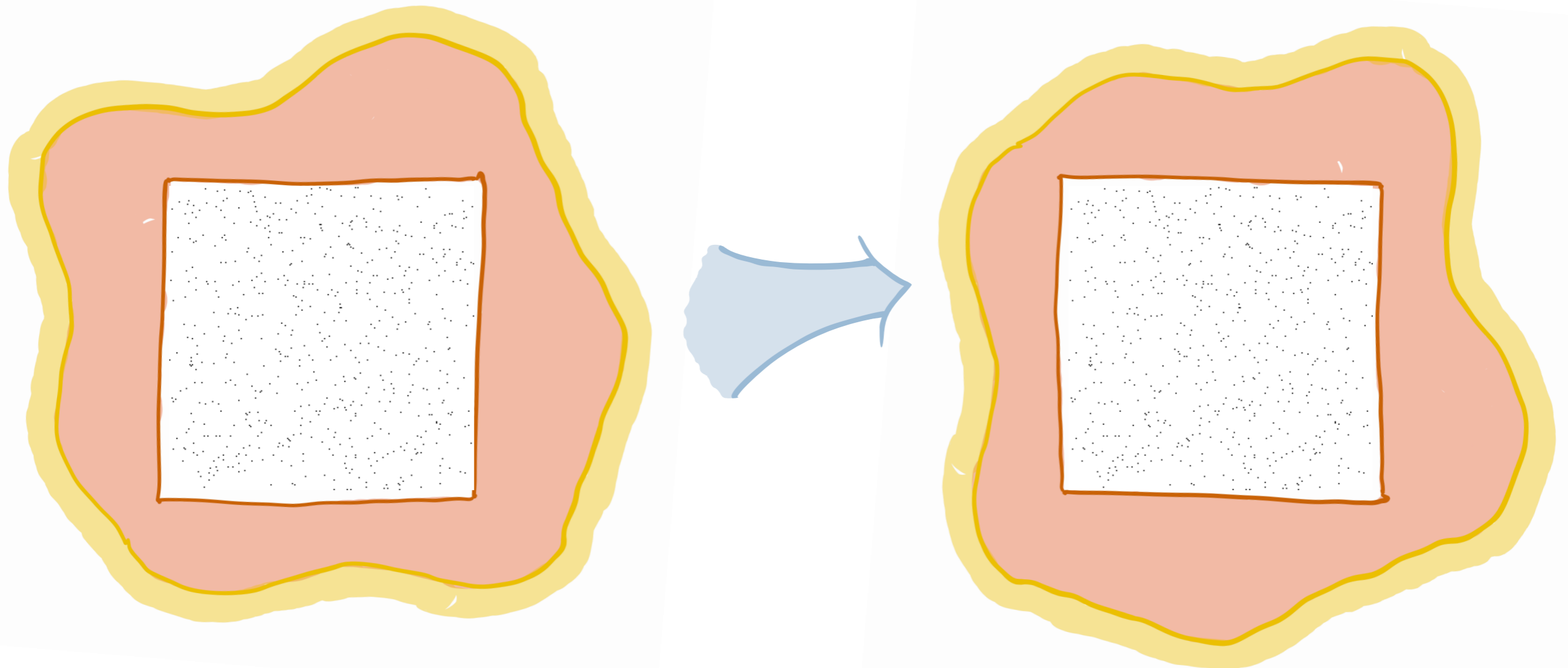


Symmetry

Simulating Irreversibility

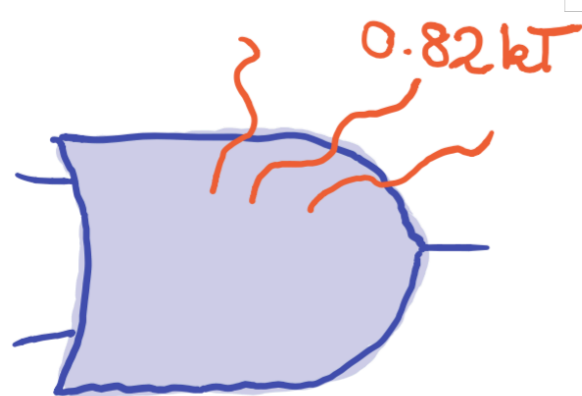
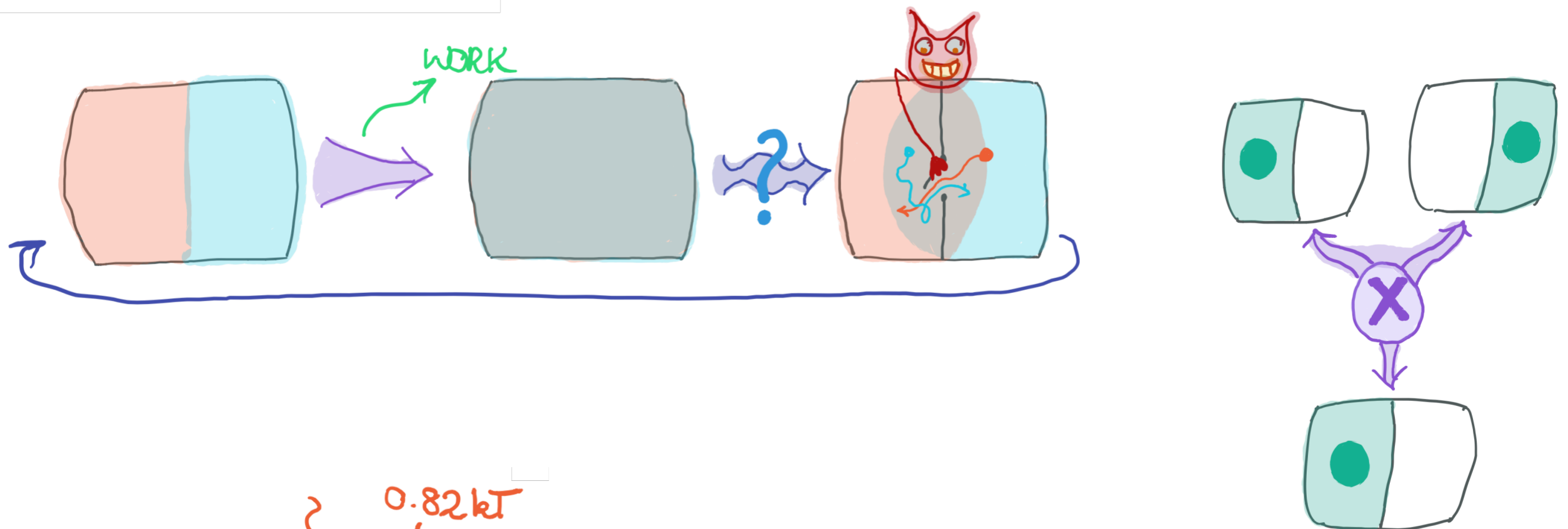


Simulating Irreversibility



Maxwell's Dæmon

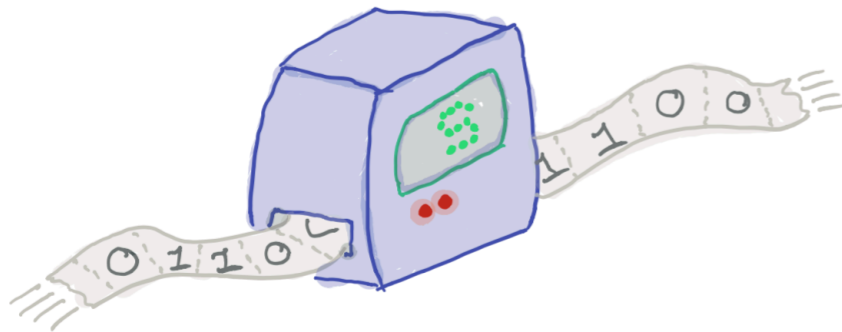
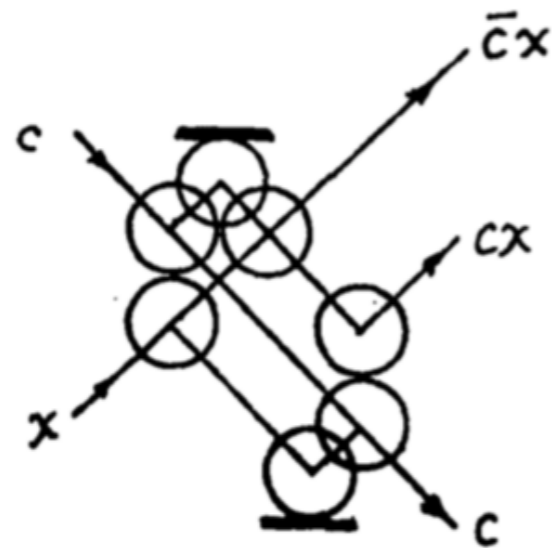
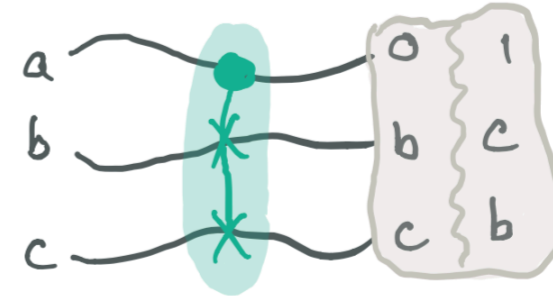
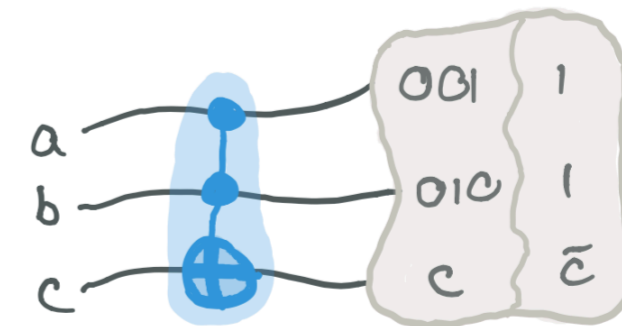
Connecting Thermodynamics to Information Theory



$$I = - \sum p_i \log p_i$$

$$\Delta q_h \geq k_B T \Delta I$$

Computing Reversibly

Bennett's RTM²CSWAP³ / Fredkin GateSwitch Gate³CCNOT¹ / Toffoli Gate

¹Rolf Landauer — 1961 — 'Irreversibility and heat generation in the computing process' — IBM J. Res. Dev.

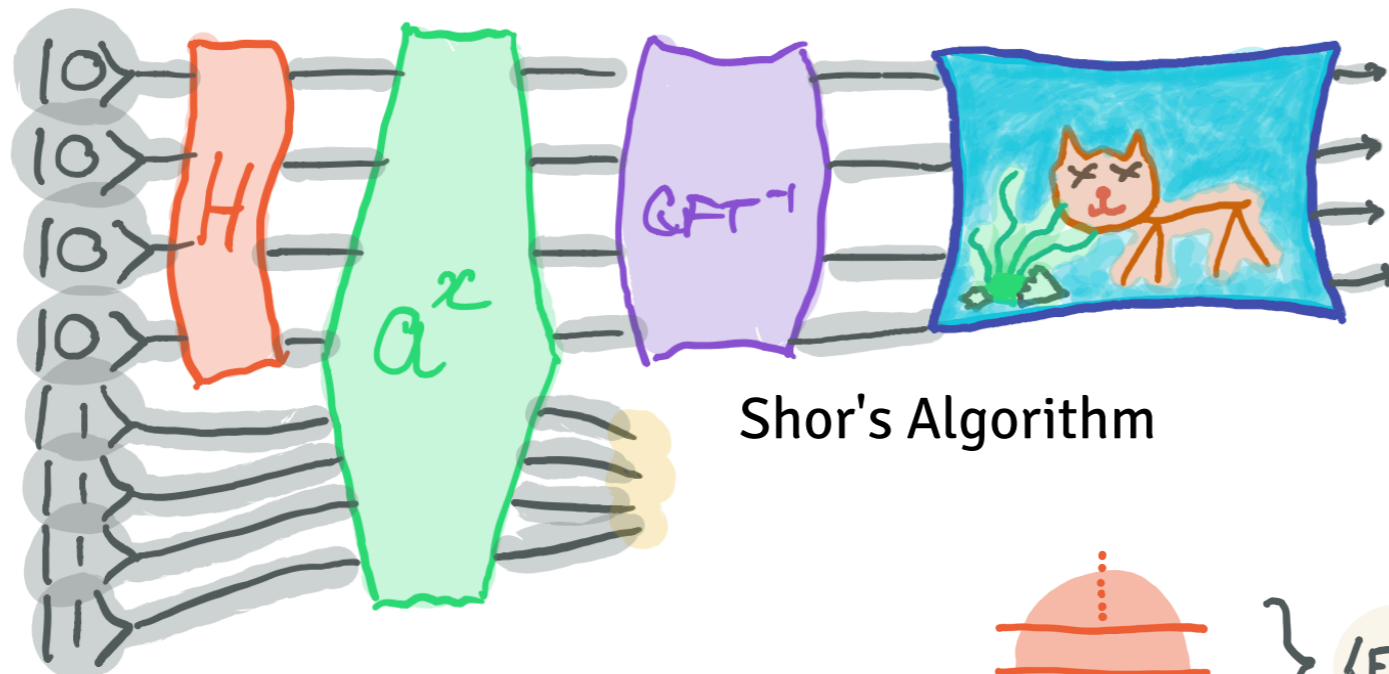
²Charles H Bennett — 1973 — 'Logical Reversibility of Computation' — IBM J. Res. Dev.

³Edward Fredkin and Tommaso Toffoli — 1981 — 'Conservative Logic' — Collision-Based Computing

“Even if classical balls could be shot with perfect accuracy into a perfect apparatus, fluctuating tidal forces from turbulence in the atmospheres of nearby stars would be enough to randomise their motion within a few hundred collisions. Needless to say, the trajectory would be spoiled much sooner if stronger nearby noise sources (e.g., thermal radiation and conduction) were not eliminated.”

– *Charles Bennett*

Quantum Computers



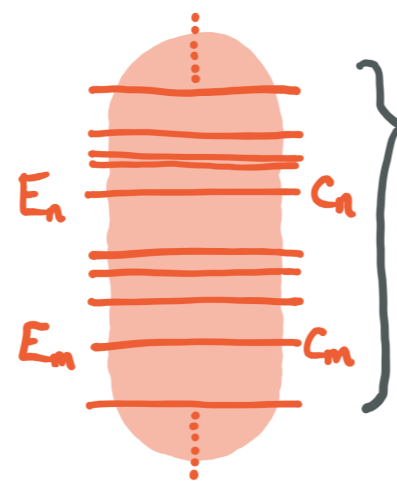
Shor's Algorithm

$$i\hbar\partial_t |\psi(t)\rangle = \mathbf{H} |\psi(t)\rangle$$

$$|\psi(n \delta t)\rangle = \mathbf{U}^n |\psi(0)\rangle$$

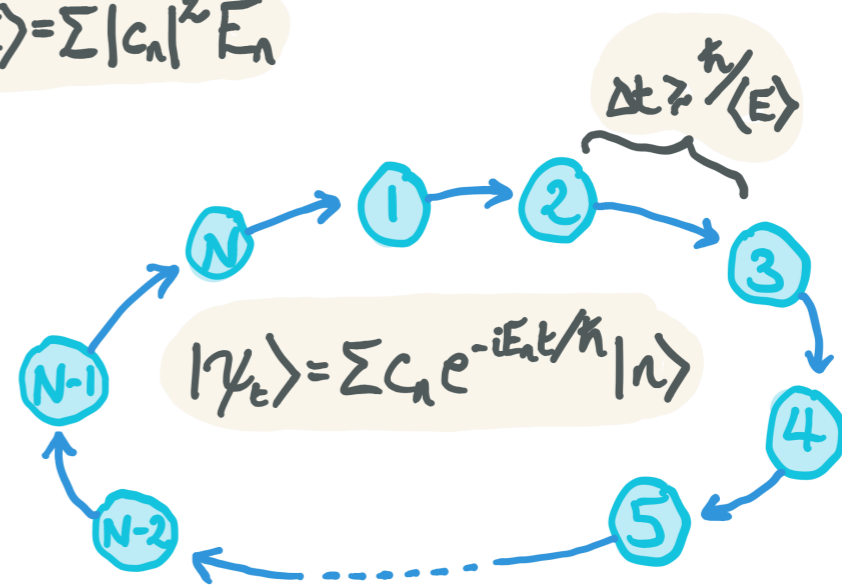
$$\mathbf{U} = e^{-i\mathbf{H}\delta t/\hbar}$$

$$\mathbf{U}^{-1} \equiv \mathbf{U}^*$$



$$\langle E \rangle = \sum |c_n|^2 E_n$$

$$\langle m | n \rangle = \delta_{mn}$$



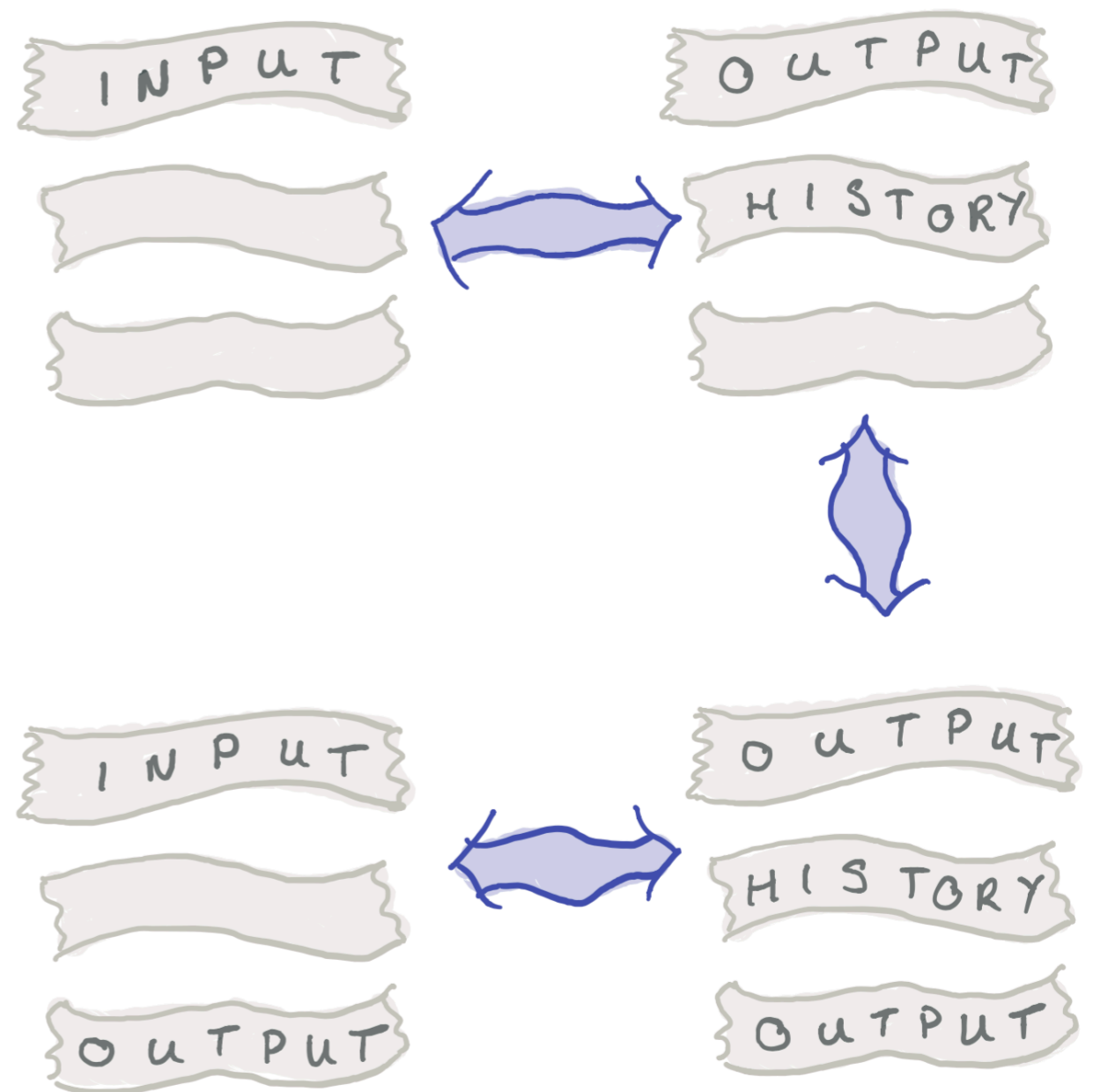
Garbage



Garbage

$$f : x \mapsto f(x)$$

$$\tilde{f} : x \leftrightarrow (x, f(x))$$



Pebbling

TABLE 2
Reversible simulation in time $O(T^{\log 3/\log 2})$ and space $O(S \cdot \log T)$.

Stage	Action	Checkpoints in storage (0 = initial ID, checkpoint $j = (jm)$ th step ID)								
0	Start	0								
1	Do segment 1	0	1							
2	Do segment 2	0	1	2						
3	Undo segment 1	0		2						
4	Do segment 3	0		2	3					
5	Do segment 4	0		2	3	4				
6	Undo segment 3	0		2		4				
7	Do segment 1	0	1	2		4				
8	Undo segment 2	0	1			4				
9	Undo segment 1	0				4				
10	Do segment 5	0				4	5			
11	Do segment 6	0				4	5	6		
12	Undo segment 5	0				4		6		
13	Do segment 7	0				4		6	7	
14	Do segment 8	0				4		6	7	8
15	Undo segment 7	0				4		6		8
16	Do segment 5	0				4	5	6		8
17	Undo segment 6	0				4	5			8
18	Undo segment 5	0				4				8
19	Do segment 1	0	1			4				8
20	Do segment 2	0	1	2		4				8
21	Undo segment 1	0		2		4				8
22	Do segment 3	0		2	3	4				8
23	Undo segment 4	0		2	3					8
24	Undo segment 3	0		2						8
25	Do segment 1	0	1	2						8
26	Undo segment 2	0	1							8
27	Undo segment 1	0								8

Thinking Reversibly

- **Bennett's algorithms**
 - efficient embedding of irreversibility
 - not easily composable
 - injective rather than bijective



Thinking Reversibly

- True reversible programming: make use of bijections

$$+ : (a, b) \mapsto (a, b, a + b)$$



$$+_1 : (a, b) \leftrightarrow (a + b, b)$$

$$+_2 : (a, b) \leftrightarrow (a + b, a - b)$$

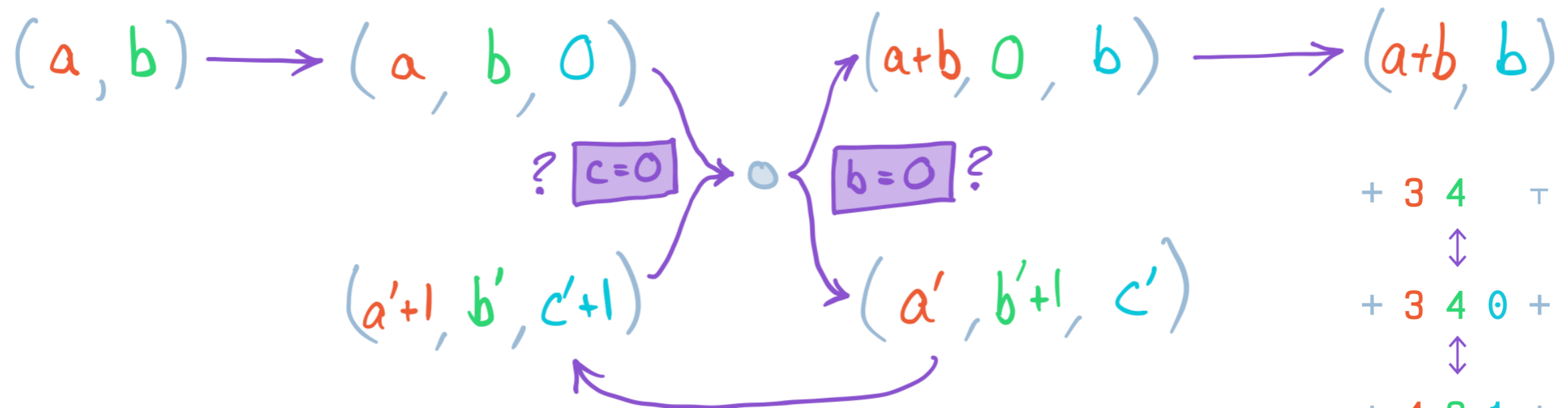


- By appropriately exploiting information in the output, can reduce or even eliminate the 'garbage' output
- The remaining garbage, if purposefully constructed, is often found to be useful for further computation
- For example, it turns out that $+_1$ and some reversible looping is sufficient for

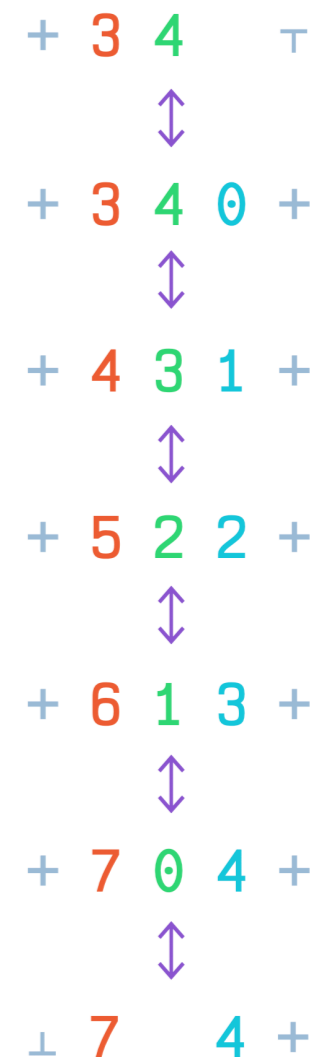
$$\text{square} : n \leftrightarrow n^2$$

Example 1

Addition/Subtraction



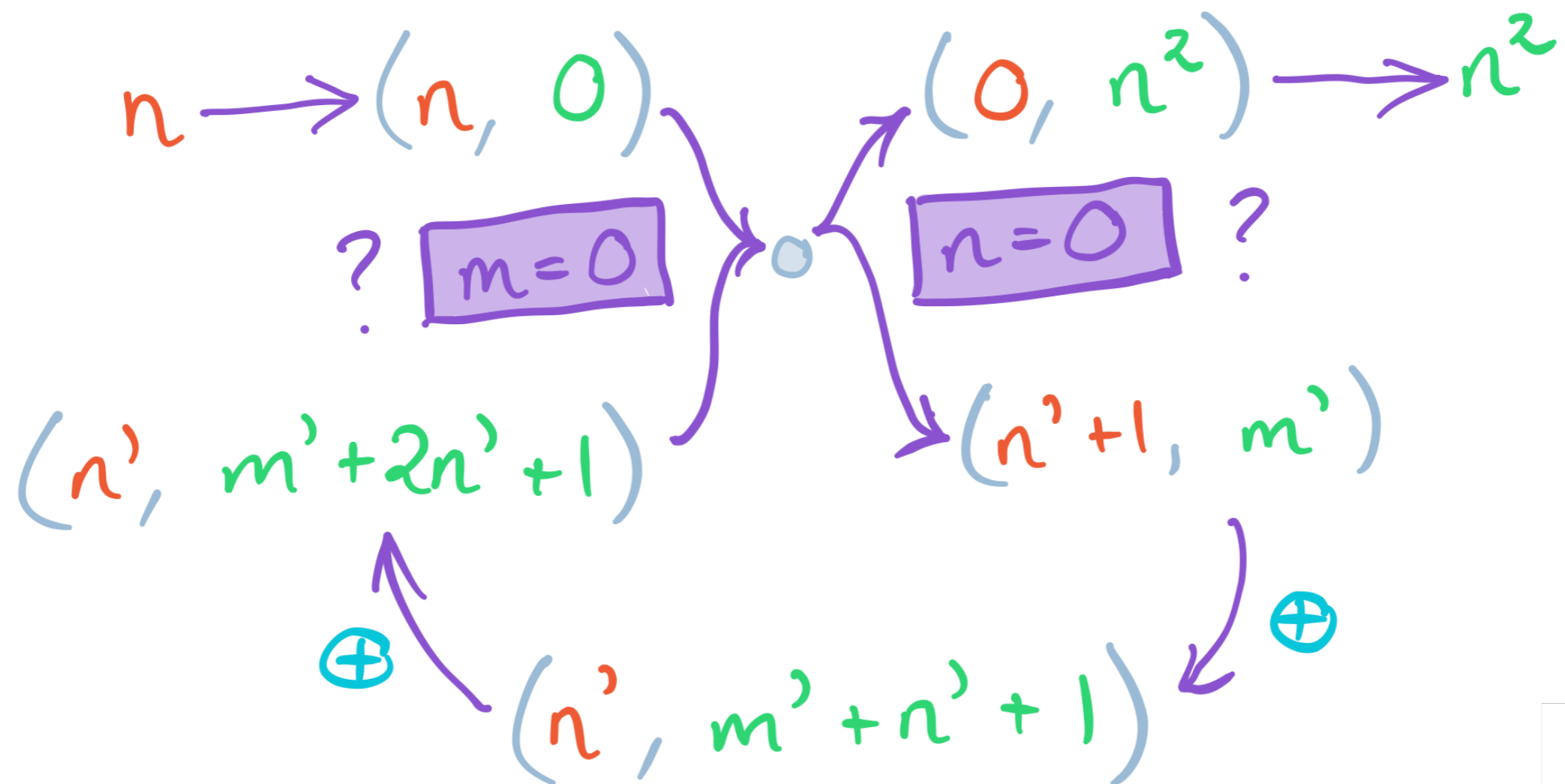
- Reversible analogue of Peano definition
- Straightforward extension to Integers
- Can also implement for Rationals and Reals



Example 2

Square/Square Root

$$n^2 = \sum_{k=0}^{n-1} (2k + 1)$$



$$i\hbar \frac{\partial}{\partial t} |\psi\rangle = \underbrace{\left(-\frac{\hbar^2}{2m} \frac{\partial^2}{\partial x^2} + V(x) \right)}_{\mathbf{H}} |\psi\rangle$$

Example 3

Schrödinger's Equation

$$i\hbar \partial_t |\psi(t)\rangle = \mathbf{H} |\psi(t)\rangle$$

$$|\psi(n \delta t)\rangle = \mathbf{U}^n |\psi(0)\rangle$$

$$\mathbf{U} = e^{-i\mathbf{H}\delta t/\hbar}$$

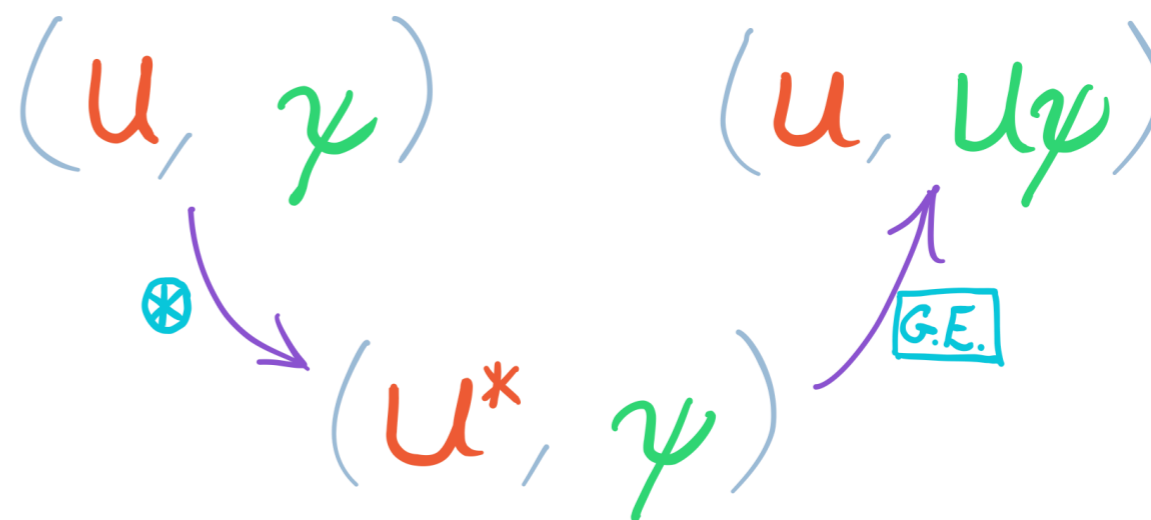
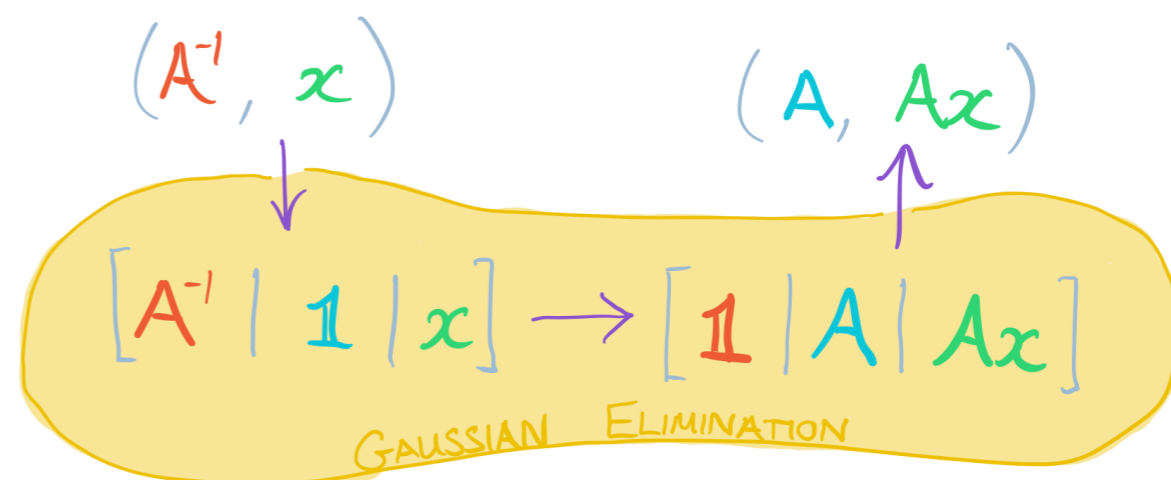
$$\mathbf{U}^{-1} \equiv \mathbf{U}^*$$

```
;; This subroutine changes a point in the real wave DEST
;; according to the curvature in the corresponding
;; neighborhood in the real wave SRC, and the potential at
;; the given point.
(defun pfunc (dest src i alphas epsilon)
  ((dest _ i) += ((alphas _ i) */ (src _ i)))
  ((dest _ i) -= (epsilon */ (src _ ((i + 1) & 127))))
  ((dest _ i) -= (epsilon */ (src _ ((i - 1) & 127))))))

;; Take turns updating the two components of the wave in a
;; way such that they will chase each other around in
;; (higher-dimensional) circles.
(defun schstep (psiR psiI alphas epsilon)
  ;; psiR += f(psiI)
  (for i = 0 to 127
    ;; psiR[i] += pfunc(psiI,i)
    (call pfunc psiR psiI i alphas epsilon))
  ;; psiI -= f(psiR)
  (for i = 0 to 127
    ;; psiI[i] -= pfunc(psiR,i)
    (rcall pfunc psiI psiR i alphas epsilon)))

;; Print the current wave to the output stream.
(defun printwave (wave)
  (for i = 0 to 127
    (printword (wave _ i)))
  (println))

;; Main program, goes by the name of SCHROED.
(defun schroed
  (for i = 1 to 1000 ;Time for electron to fall to well bottom.
    (call schstep psiR psiI alphas epsilon)
    ;; Print both wave components.
    (call printwave psiR)
    (call printwave psiI)))
```



Reversible Languages

```
(defsub mult (m1 m2 prod)
  ;; Use grade-school algorithm:
  (for pos = 0 to 31
    (if (m1 & (1 << pos)) then
      (prod += (m2 << pos))))))
```

R¹

```
(defun fact (n) (assert (and (integerp n) (> n 0)))
  (if (onep n) #'onep n (* n (fact (1- n)))))
```

Ψ-LISP²

```
procedure fib
  if n=0 then x1 += 1
              x2 += 1
  else n -= 1
        call fib
        x1 += x2
        x1 <=> x2
fi x1=x2

procedure main_fwd
  n += 4
  call fib

procedure main_bwd
  x1 += 5
  x2 += 8
uncall fib
```

Janus³

```
type Nat4 = Bool * Bool * Bool * Bool
```

```
add1 :: Nat4 ↔ Nat4 :: sub1
| (a, b, c, False)      ↔ (a, b, c, True)
| (a, b, False, True)   ↔ (a, b, True, False)
| (a, False, True, True) ↔ (a, True, False, False)
| (False, True, True, True) ↔ (True, False, False, False)
| (True, True, True, True) ↔ (False, False, False, False)
```

Theseus⁴

¹Michael P Frank — 1997 — ‘The R Programming Language and Compile’ — MIT Rev. Comp. Proj. Memo

²Henry G Baker — 1992 — ‘NREVERSAL of Fortune—the Thermodynamics of Garbage Collection’ — Intl. Workshop on Memory Management

³Tetsuo Yokoyama and Robert Glück — 2007 — ‘A reversible programming language and its invertible self-interpreter.’ — Partial evaluation and semantics-based program manipulation.

⁴Roshan P James and Amr Sabry — 2014 — ‘Theseus: a high-level language for reversible computation.’ — Reversible Computation

Reversible Languages

```

str(s|count) {
  s [
    s temp s temp s temp s temp
    s temp s temp s temp s temp
    count str(s|count)len | count
    temp s temp s temp s temp s
    temp s temp s temp s temp s
  ] s
} (s|count)len

revloop(in|count|out) {
  count [
    in temp in temp in temp in temp in
    temp in temp in temp in temp in
    out temp out temp out temp out temp out
    temp out temp out temp out temp out
    revloop(in|count|out)poolver
  ] count
} (in|count|out)poolver

reverse(in|out) {
  str(in|count)len
  revloop(in|count|out)poolver
  nel(count|out)rts
} (in|out)esrever

(ent|in) { reverse(in|out)esrever } (out|ent)

```

Kayak⁵

swap-fl1 swap-fl2 : {a b c : U} → PLUS a (PLUS b c) ↔ PLUS c (PLUS b a)
 swap-fl1 = $\text{assocl}_+ \odot \text{swap}_+ \odot (\text{id} \leftrightarrow \oplus \text{swap}_+)$

swap-fl2 = $(\text{id} \leftrightarrow \oplus \text{swap}_+) \odot$
 $\text{assocl}_+ \odot$
 $(\text{swap}_+ \oplus \text{id} \leftrightarrow) \odot$
 $\text{assocr}_+ \odot$
 $(\text{id} \leftrightarrow \oplus \text{swap}_+)$

Π⁶

-- Peano definition of *natural numbers*

```

data Z;
data S n;

```

-- reversible Peano addition

```

+ a b _ = + a b Z +;
+ (S a) b c + = + a (S b) (S c) +;
+ Z a+b a + = _ a a+b +;

```

-- termination conditions

```

! + a b _;
! _ a a+b +;

```

```

+ 3 4 _
  ↓
+ 3 0 4 +
  ↓
+ 4 1 3 +
  ↓
+ 5 2 2 +
  ↓
+ 6 3 1 +
  ↓
+ 7 4 0 +
  ↓
_ 7 4 +

```

alethe⁷

⁵Ben Rudiak-Gould — 2002 — Esoteric Programming Language Awards

⁶Jacques Carette, Roshan P. James, and Amr Sabry — 2018 — 'Embracing the laws of physics: Three reversible models of computation.' — arXiv

⁷William Earley — 2019 — DNA25 Poster; Paper in Preparation

Quantum Languages

```

procedure grover(int n) {
  int l=floor(log(n,2))+1;    // no. of qubits
  int m=ceil(pi/8*sqrt(2^l)); // no. of iterations
  int x;
  int i;
  qureg q[1];
  qureg f[1];

  {
    reset;
    Mix(q);                // prepare superposition
    for i= 1 to m {        // main loop
      query(q,f,n);        // calculate C(q)
      CPhase(pi,f);        // negate |n>
      !query(q,f,n);       // undo C(q)
      diffuse(q);          // diffusion operator
    }
    measure q,x;          // measurement
    print "measured",x;
  } until x==n;
}

```

QCL

Solve a circuit-satisfiability problem.

```

!include <gates>

!use_macro not1 not_x4
not_x4.$A = x3
not_x4.$Y = $x4

!use_macro or2 or_x5
or_x5.$A = x1
or_x5.$B = x2
or_x5.$Y = $x5

```

QMASM

import Quipper

```

spos :: Bool -> Circ Qubit
spos b = do q <- qinit b
          r <- hadamard q
          return r

```

Quipper

$$\text{toff} : \mathbb{Q}_2 \multimap \mathbb{Q}_2 \multimap \mathbb{Q}_2 \multimap \mathbb{Q}_2 \otimes (\mathbb{Q}_2 \otimes \mathbb{Q}_2)$$

$$\text{toff } c \ x \ y = \text{if}^\circ c$$

$$\text{then } (\text{qtrue}, \text{cnot } x \ y)$$

$$\text{else } (\text{qfalse}, (x, y))$$

QML

Q#

```

def solve[n:!N](bits:!B^n){
  // prepare superposition between 0 and 1
  x:=H(0:B);
  // prepare superposition between bits and 0
  qs := if x then bits else (0:int[n]) as B^n;
  // uncompute x
  forget(x=qs[0]); // valid because `bits[0]==1`
  return qs;
}

```

Silq

qCGL

LIQUiD

Q|SI>

QPL

Q

QFC

Reversible vs Irreversible

- Equipotent
- Fundamentally same resource usage
 - Reversible makes erasure explicit
 - Can reduce net resource usage
- Reversibility *encourages* a more careful and efficient approach to resource usage

Hybrid

- What might a reversibility exploitation story look like?
 - Reversible processor with entropy-dissipation co-processor
 - Expose irreversible abstraction layer over reversible operations (or reversibility aware compiler to insert erasure operations)
 - Gradual software transition
- Intermediate benefit: transition from actively dissipative transistors to low energy components

Linear Types

- Linear logic is a superset of Quantum logic
 - Capable of describing reversible computation
 - Not strictly reversible though, e.g.
 $£1 \multimap (\text{sweets \& crisps \& drink})$
- Gateway drug to reversible programming
 - e.g. new Haskell extension to specify that an argument *must* be consumed (forbids *implicit* erasure)

Research Timeline

Theory

- Information–Entropy Connection (Szilard '29, Landauer '61)
- Foundation of Reversible Computation (Bennett '73)
- Ballistic Computation (Toffoli+Fredkin '81)
- Programming Languages (Ψ -LISP: Baker '92, R: Frank '97, ...)
- Analyses of Physical Limits (Frank '99, Lloyd '02, ...)

Research Timeline

Application

- Helical Logic (Merkle '96)
- Passive Transistor Logic (de Vos et al '99)
- Quantum Dot Cellular Automata (Lent et al '01-'03)
- nSQUID Josephson-Junction Circuits (Semenov et al '03)
- Asynchronous Ballistic Fluxon Logic (Frank et al '17)

...along with research into effective synthesis of the relevant circuits

Engines of Parsimony

- I. How fast can any computer run, given some spacetime region and power supply and taking into account all areas of physics?
- II. What happens when we try to communicate between/synchronise reversible computers?
- III. What happens when we try to share some common resource between asynchronous reversible computers?

Thank you!



**UNIVERSITY OF
CAMBRIDGE**

EPSRC

Engineering and Physical Sciences
Research Council



**Department of Applied Mathematics
and Theoretical Physics (DAMTP)**